Test Method for Process Deadlock Based on Graph Grammars

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Abstract
This paper proposes a test method for process deadlock based on graph grammars through constructing process resource diagram. Through using the construction rules, it can construct and judge the validity of the process resource diagram. Through using the test rules, it can test if there is the deadlock in the process. The method is a graphical approach; it is simple and intuitive with strong operability.

Keywords: process, deadlock, test, graph grammars

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S: E → V, T: E → V are two functions that gives the start node and the end node of each edges of a graph G;

**Definition 2:** A production of a graph grammar is a rule likes gl : = gr. gl and gr are two graphs, called as the left and the right of the production separately.

**Definition 3:** Host graph, expressed as ghost, is a graph that will be transformed using productions. In additional, a subgraph of a graph that is isomorphic with gl is expressed as ghost, which will be replaced as a graph that isomorphic with gr when a graph is transformed.

**Definition 4:** From ghost, remove all the nodes and the edges of a glhost and the edges which endpoint on glhost will get a graph, called as Residual Graph, marked as gresidual.

**Definition 5:** ghost is a graph, which is isomorphic with gr, and replaces ghhost of a host graph.

The productions is the basis of the transformation for a host graph, but in order to complete the transformation, only production is not enough, the corresponding rules also should explain how to embed ghghost into gresidual, the rules called embedding rules.

**Definition 6:** A graph grammar gg is a triple (A, P, E), among them, the A is the initial graph of the graph grammar, the P is the set of the productions, the E is the embedding rules.

**Definition 7:** Set gg = (A, P, E) is a graph grammar, L (gg) = {G| A => G ∧ there is not non-terminal in G} is a language set of a graph grammar.

For a graph grammar, there are two operations: one is derivation, the other is parsing. The prior means that from the host graph to find out a subgraph that isomorphic with the left graph of the production and replaced with a graph that isomorphic with the right graph of the production; the latter means that from a graph to find out a subgraph that isomorphic with the right graph of the production and replaced with a graph that isomorphic with the right graph of the production. The graphs that derived from the host graph called as the language of the graph grammar. Through parsing, if a graph can arrive to the initial character, the graph is called as a language of the graph grammar.

### 3. Construction for Process Resource Diagram Based on Graph Grammar

**Definition 1:** A process resource diagram (PRD) can be represented as PRD=(P, R, Er, Ed), among them:
- P is a set of processes;
- R is a set of resources;
- Er is a set of edges according to a process requests for resources;
- Ed is a set of edges according to a resource assigned to a process.

For example, Figure 1 is a PRD,

![Figure 1. A PRD](image)

Among them, P= {p1, p2}, R= {R1, R2}, Er= {<p1, R1>, <p1, R2>, <p2, R1>, <p2, R2>}, Ed= {<R1, p1>, <R2, p2>},

- `<pi, Rj>` presents a process pi applies a resource Rj;
- `<Rj, pi>` presents a resource Rj has been allocated to a process pi.
Definition 2: A reasonable process resource diagrams is a PRD in that the processes have not been deadlocked in a state. A reasonable process resource diagram should satisfy two conditions:

The allocation for a resources Rj can not exceed the total;
The sum of the number of any process pi applies for a resource Rj and the number of the resource Rj allocate to the processes is no more than the total of the resource Rj.

Definition 3: A complete schematic diagram is a PRD that contains only an isolated node.

Based on the basic principle of graph grammars, through the rules and using the derivation of the graph grammar it can construct the process resource diagram. Design rules are as Figure 2.

Figure 2. A Set of the Rule to Construct the Process Resource Diagram

Among them, $P$ means a process; $R:m$ means a kind of resource, m is the number of the resource; $r$ is a non-terminal node that of a resource.

Rule 1 means that it begin to draw a resource, that is to say to draw a process resource diagram;
Rule 2 means starting from a non-terminal node of a resource and bring in a terminal node;
Rule 3 means that all the non-terminal nodes of the resources changed into the terminal nodes and the resources had been brought end;
Rule 4 bring in a process and the process puts forward an application to a resource;
Rule 5 means that an existing process puts forward an application to a resource;
Rule 6 distributes a resource to a process;
Rule 7 means that an application that a process puts forward to a resource been changed to the allocation that a resource to a process.

Figure 3 is the construction process of the Figure 1,

\[ \lambda \xrightarrow{(1)} r \xrightarrow{(2)} R1:2 \quad r \xrightarrow{(3)} R1:2 \quad R2:3 \]

\[ \begin{array}{c}
P1 \\
P2 \\
R1:2 \\
R2:3 \\
\end{array} \quad \begin{array}{c}
P1 \\
P2 \\
R1:2 \\
R2:3 \\
\end{array} \quad \begin{array}{c}
P1 \\
P2 \\
R1:2 \\
R2:3 \\
\end{array} \]

Figure 3. The Construction Process of the Figure 1

4. Test for Process Deadlock Based on Graph Grammar

**Theorem**: The system is a deadlock state, if and only if its process resource diagram can not change to a complete schematic diagram.

Based on the theorem, we can judge if it have a deadlock in a system.
Step 1: In the process resource diagram to look for if it have a process node that the sum of the application edges and the allocation edges less than the total of the resource. If it is yes then it turn to step2, else it turn to step3;
Step 2: For the process node that satisfy the condition, we can modify the application edges to the allocation edges, and this makes the process node be an isolated node; Turn to step 1;
Step 3: If it is a complete schematic diagram, the system does not exist the deadlock.

Based on the basic principle of graph grammars, through the rules and using the parsing of the graph grammar it can construct the process of the test for process deadlock. Design rules are as Figure 4. Among them:

- Rule 1 puts an application edge change to an allocation edge for a process node;
- Rule 2 deletes all the allocation edges of a resource allocate to a process;
- Rule 3 put two isolated nodes parse to a isolated node;
- Rule 4 put a terminal node resource parse to a non-terminal node resource;
Rule 5 put a terminal node resource and a non-terminal node resource parse to a non-terminal node resource;
Rule 6 put an isolated resource and an isolated process parse to the initial character.
From the process of the test for process deadlock of the Figure 1, we can see that the Figure 1 can be parsed to the initial graph: $\lambda$, so the Figure 1 does not exist the deadlock.

Figure 4. A set of the Rule to Construct the Test for Process Deadlock

Figure 5. The Test for Process Deadlock of the Figure 1

5. Summarize and Look Forward
This paper proposes a test method for process deadlock based on graph grammars through constructing process resource diagram. Through using the construction rules, it can construct and judge the validity of the process resource diagram. Through using the test rules, it can test if there is the deadlock in the process. The method is a graphical approach; it is simple and intuitive with strong operability.
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